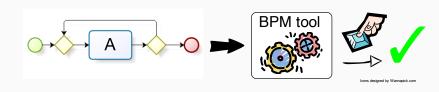
# Add data into business process verification

(ab)using planning tools for BPM

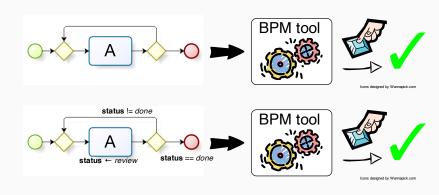
**Sergio Tessaris** joint work with Riccardo De Masellis, Chiara Di Francescomarino, Chiara Ghidini KRDB SOS-2020 11/9/2020

online slides available on stessaris.pages.scientificnet.org/talks/sos2020

### The problem with current tools



## The problem with current tools



- Theory is there
  - why not available?
  - what can we do to bridge the gap?
- Scalability problem?

#### Our research

- Lots of *implementable* frameworks and methods
- Do available tools scale up?
- How do we verify that?

## Reachability verification

- Interesting tasks can be reduced to reachability; e.g.
  - (proper) termination
  - dead transitions
  - trace repair/completion
  - ...

### E.g. Trace completion

- Assume model
- Given partial log
  - sequence of events + data updates
  - empty, partial, or complete
- Find complete sequence compatible with log

#### **Classical planning**

- Essentially reachability verification
- Plans/strategy as bonus
  - e.g. answer what to do next questions
- Strong community interested on scalability
  - International Planning Competition running since 1998

### Planning and workflows

- Planning for workflows
  - On the Disruptive Effectiveness of Automated Planning for LTLf-Based Trace Alignment De Giacomo et al. AAAI 2017
  - Automated Planning for Business Process Management Marrella Journal on Data Semantics, June 2019
- Workflows for planning
  - Planning via Petri Net Unfolding Hickmott et al. IJCAI 2007

#### Focus of this work

- How we can exploit the available planning tools
- Which tools are best suited for workflow analysis
- Evaluate scalability

#### Putting tools to the test

- Focus on reachability; i.e. automated planning
- Select industrial-strength tools:
  - Answer Sets Programming: Clingo
  - Classical Planning: Fast Downward
  - Model checking: nuXmv
- Build a common ground among the tools
- Select appropriate experiments

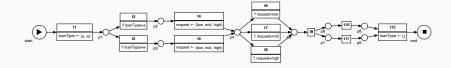
#### Data-aware Workflows: which one?

- Several proposed frameworks
- We selected a simple one
  - i.e. close to Classical Planning
- It's an initial step

## DAta Workflow Nets (DAWNets)

- Workflow Nets
  - connected Petri Nets, with start and stop places
- Variables
  - domain
  - possibly unassigned
  - transitions assign values
- Transition guards
- based on Soundness verification for conceptual workflow nets with data Sidorova et al. Information Systems 2011/11

# **DAWNet** example



#### **DAWNets semantics**

- Extends PN semantics
  - State: marking + vars
- Valid firing  $t: s \leadsto s'$ 
  - s: token in each input place of t
  - s: guard is satisfied
  - s': tokens from in to out places
  - s': variables updated
- Case: sequence of valid firing

#### Restrictions

- Bounded networks (safeness)
  - correct algorithms to check it
- Finite domains

### Different paradigms one task

- Different tools uses different languages
  - several ad hoc encodings in literature
- Common denominator: Labelled Transition Systems

### **DAWNets reachability as LTS**

- Labels: transition names
- States:  $(M, \nu)$  marking + variable assignment
- Initial state: token in start + unassigned variables
- Transition relation: (s, t, s') based on firing  $t: s \rightsquigarrow s'$
- Goal states satisfying required properties
  - e.g. proper termination

#### **PDDL Planners**

- Reachability in LTS is a planning problem
- Actions schemata
  - pre/post conditions
- Initial conditions
- Final conditions
- Built-in frame axiom
- Operational semantics

#### PDDL in practice

- Planners are optimised for subsets of the language
- Fast Downward
  - grounding!
- Several heuristics, some depending on PDDL subset
- E.g. no object fluents
  - only boolean predicates
- Places: constants + active predicate
- Transitions: actions
- Variables: unary predicates

### Planning using ASP

- Fluents
- Causation rules to define the LTS
  - head depends on both previous and current states

```
F if G ifcons H after M.
t:F:-t:G, not not t:H, (t-1):M.
```

- Variables, ASP style strong negation
  - grounding!
- Valid states are stable models wrt the rules
- Compact encoding
- Native domain constraints
- Inertia is not builtin

### **ASP** planning in practice

- Language is not standardised
- Coala (based on Clingo)
  - Not optimised
- Places: unary predicate
- Transitions: actions
- Variables: unary predicates

### Model checking

- Tools are based on TL over infinite traces (LTL)
  - looping on final states
- Variables over arbitrary domains or booleans
- TS defined using formulae over current and previous states
- Native constraints over domain
- Inertia is not builtin
  - no NMR to help with that

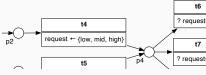
## Model checking in practice

- nuXmv
- Places: boolean variables
- Transitions: variable over transition names
- Variables

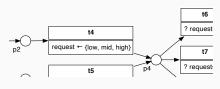
## **Encodings**

- More details in my early talk
- Leveraging trace equivalence
  - formally proven for each encoding
- Transferable models

# PDDL encoding (Domain)



```
(:constants
  p2 p3 p1 p6 p7 p4 p5 p8 p9 start end - place
  high s low w mid - active_domain
(:predicates
  :: Places
  (p_enabled ?p - place )
  (p_terminal ?p - place )
  :: Variables
  (request ?v - active_domain )
  ;; Domains
  (t4_request_domain ?v - active_domain)
```



```
(:action t4
  :parameters ( ?request - active_domain )
  :precondition
    (and (p_enabled p2) (t4_request_domain ?request))
  :effect
    (and
          (p_enabled p4)
          (not (p_enabled p2))
          (forall (?v - active_domain) (not (request ?v)))
          (request ?request)))
```

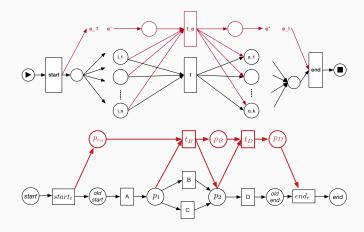
### PDDL encoding (Problem)

```
(:init
    (p_enabled start)
    (p_terminal end)
    (t4_request_domain low)
    (t4_request_domain mid)
    (t4_request_domain high)
(:goal
  (and
    (p_enabled end)
    (forall (?p - place)
      (or (p_terminal ?p) (not (p_enabled ?p))))))
```

## Which experiments?

- Difficult to design general reachability experiments
- Focus on *Trace Completion*
- Several parameters via traces:
  - completeness degree
  - compliance
  - size

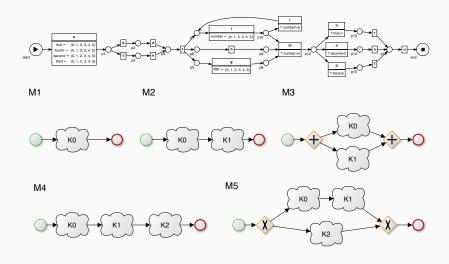
# Trace Completion as Reachability



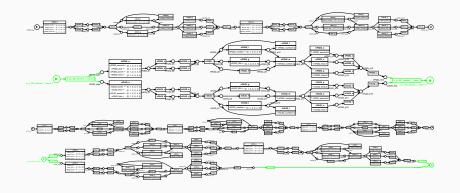
#### Synthetic models

- Synthetic base model
- Combination of copies of the base model
- 8 different traces per model (size 10-50+)
  - empty, 25%, 50%, 75%
  - also not compliant (4)

#### Base model



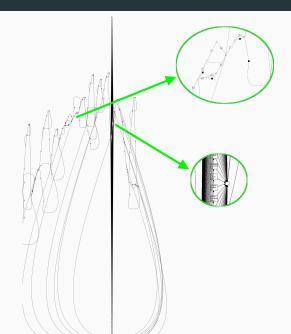
#### Models



#### Real life model

- BPI Challenge 2011 logs
  - Real life log of a Dutch academic hospital
- Model discovered using ProM Data-flow Discovery plugin
  - 355 transitions, 61 places, 710 edges, and 4 variables
- 9 random traces (size 3-500+)
  - empty, 25%, 50%, 75%
  - all compliant
- Also tested without data

# Discovered model



#### Infrastructure

It's not reproducible if it only runs on your laptop

#### Jon Zelner

- Focus on reproducibility
- Leveraging Docker, and Kubernetes cluster
- More details on Use Containers for your Experiments!

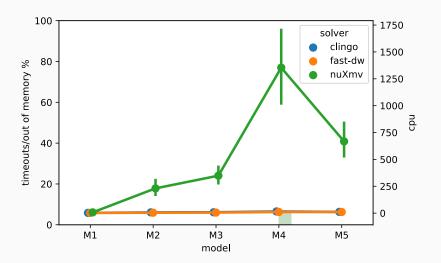
## Have we a winner?

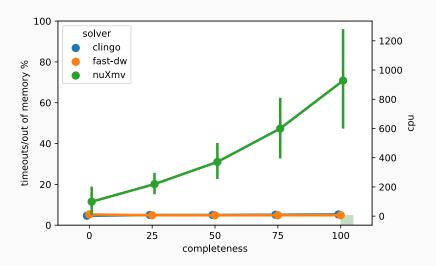
• Simple answer: **no** 

### Have we a winner?

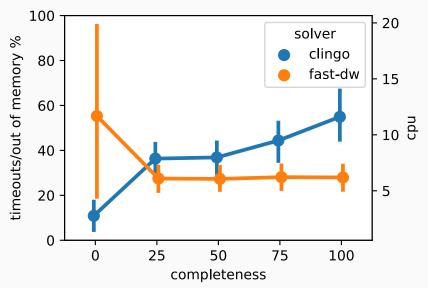
- Simple answer: no
- Actually, the picture is more complex... let's look at some pictures

# Synthetic experiments

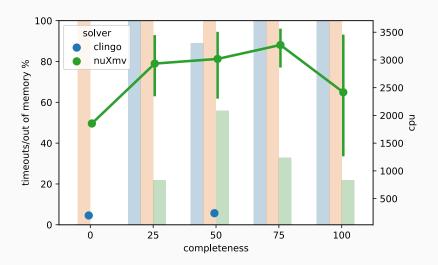




Zooming into FD/clingo

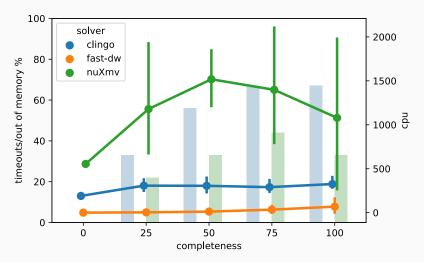


# Real life experiments



• What if we get rid of data?

• What if we get rid of data?



## Ground or not to ground

- Grounding can be expensive!
- For big problems symbolic verification can be beneficial
- Room for hybrid methods?

# **Optimising encoding**

- We didn't optimise encoding and heuristics
- Need collaboration with tool developers
  - E.g. direct encoding in SAS for FD

#### **Conclusions**

- Automated planning tools are effective
- Adding data makes the difference
- Hybrid systems can be a solution
- Now we can move to more complex languages

# Thanks!

Questions?

#### References

- Solving reachability problems on data-aware workflows
  - code available
- Use Containers for your Experiments!
- Verification of workflow nets with data